

Design of High-Speed Multiplier using Parallel Prefix Adder

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Abstract - A multiplier is one of the key hardware blocks in most digital signal processing (DSP) systems. The complexity of the circuit depends mainly on the multiplication count needed to develop the method. As the circuit complexity increases the delay in the system conjointly increases which reduces the performance. A parallel array multiplier is the solution to satisfy high speed requirements. A traditional Braun multiplier comprises of an array of 16 AND gates and 12 Full Adders of 28 Transistors. A new Braun Multiplier replaces the final stage of Full Adders with Kogge-stone adder (KSA) for quicker multiplication results. Three KSA designs are proposed using logic gates 12T XOR, 14T XOR, 22T XOR. The prototypes proposed for traditional Braun multiplier and Braun multiplier with KSA are designed in the Mentor Graphics tool for 180nm and 130nm technologies with source 1.8V and 1.3V respectively.

Key Words: Digital Signal Processing (DSP), Full Adder, Braun Multiplier, Parallel Prefix Adder, Kogge-Stone Adder (KSA).

1. INTRODUCTION

The advances in VLSI technology, each in terms of speed and size, made it attainable to implement parallel multipliers hardware. Technology development further ensures better performance characteristics and widespread use in DSP systems. It carries out operations such as accumulating the sum of multiple products much faster than a common microprocessor. The DSP architecture is designed to perform parallel operation and thus reduces the computational complexity and increase the speed required for repetitive signal processing in such applications[1]. These features are designed for higher speed and throughput in the programmable DSP. There are a large number of programmable DSPs to choose from for a given application, based on factors such as speed, throughput, arithmetic capability, precision, scale, cost and power consumption[2]. The introduction of single-chip multipliers and their incorporation into microprocessor architecture is the most important reason why commercial VLSI chips are available that are capable of DSP functions[3]. Parallel prefix adders were established as the most efficient binary add-on circuits. Their regular structure and speedy performance make them especially attractive to implement VLSI[4]. Product generation of numbers calls for a single processor cycle. Either a software-based shift and add algorithm or one

using micro-coded controllers used as popular multiplication schemes to implement the same algorithm in hardware. Both of these options require multiplication completion by several processor cycles. Kogge-stone Adder employs XOR and AND logic gates to design parallel prefix adders[5][6]. In this work, Braun Multiplier with Kogge-stone Adder is used for decreasing the delay.

2. CONVENTIONAL BRAUN MULTIPLIER

CMOS Braun multiplier consists of an array of AND gates and Full Adders. It consists of 16 AND gates and 12 Full Adders with 28 Transistors each. FA1 to FA12 are the 12 Full Adders.

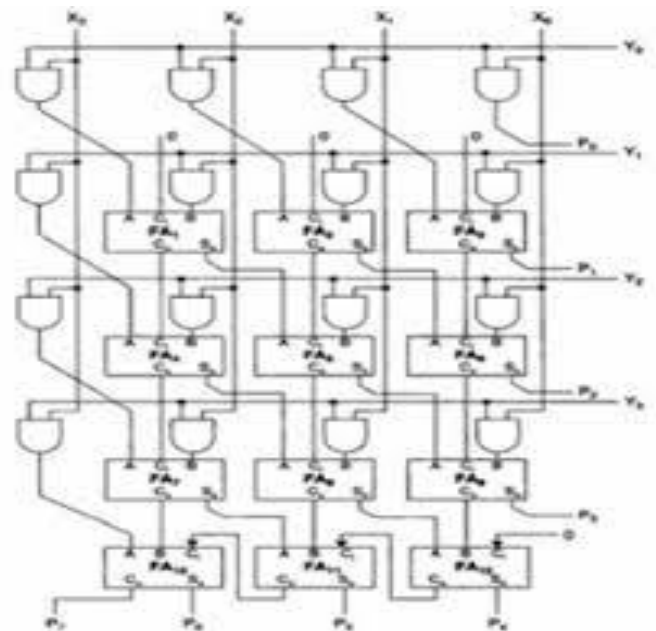


Fig -1: Circuit diagram of 4 x 4 multiplier

Consider the multiplication if two unsigned number A and B. Let number A be represented using 'm' bits ($A_{m-1} A_{m-2} \dots A_0$) in (1) and number B using 'n' bits ($B_{n-1} B_{n-2} \dots B_0$) in (2). The multiplicand A, the multiplier B, and the product P are given by (3)

$$A = \sum A_i 2^i \quad (1)$$

$$B = \sum B_i 2^i \quad (2)$$

$$P = \sum [\sum A_i B_j 2^{i+j}] \quad (3)$$

and can have a maximum of (m+n) bits.

3. SIMULATION RESULT OF CONVENTIONAL 4-BIT BRAUN MULTIPLIER

Fig. 2 represents the schematic of the 4-bit Braun multiplier with 4x4 input pins and 8 output pins. The test bench for the multiplier is shown in the fig. 3. For inputs ($X_0 - X_3, Y_0 - Y_3$) V_{pulse} is given as 1.8V and 1.3V for 180nm and 130nm technologies respectively. The output pins are given from P_0 to P_7 . Fig. 4 represents the results of the traditional Braun multiplier. The traditional Braun multiplier consists of 432 transistors and it is found that the conventional Braun multiplier has an overall delay of 535.72 psec.

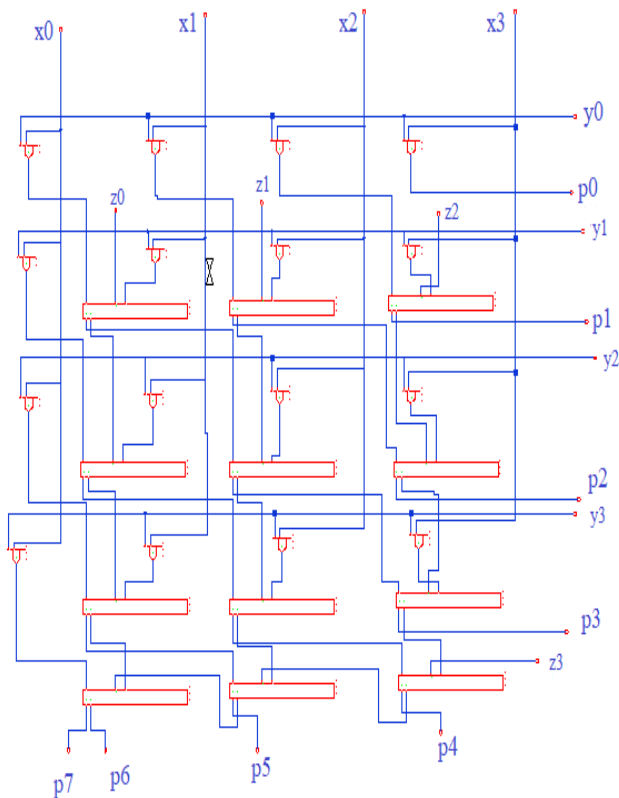


Fig -2: Schematic of Braun multiplier

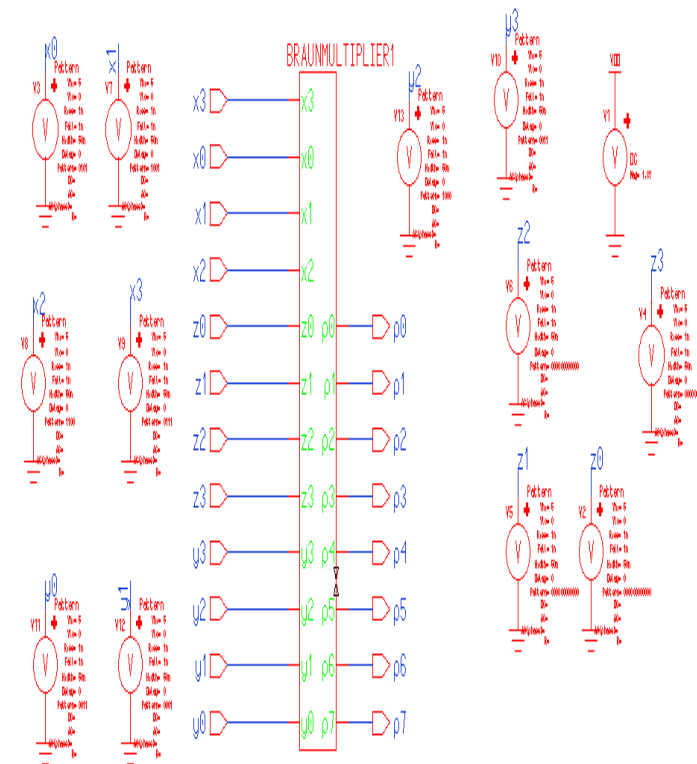


Fig -3: Test bench of Conventional Braun multiplier

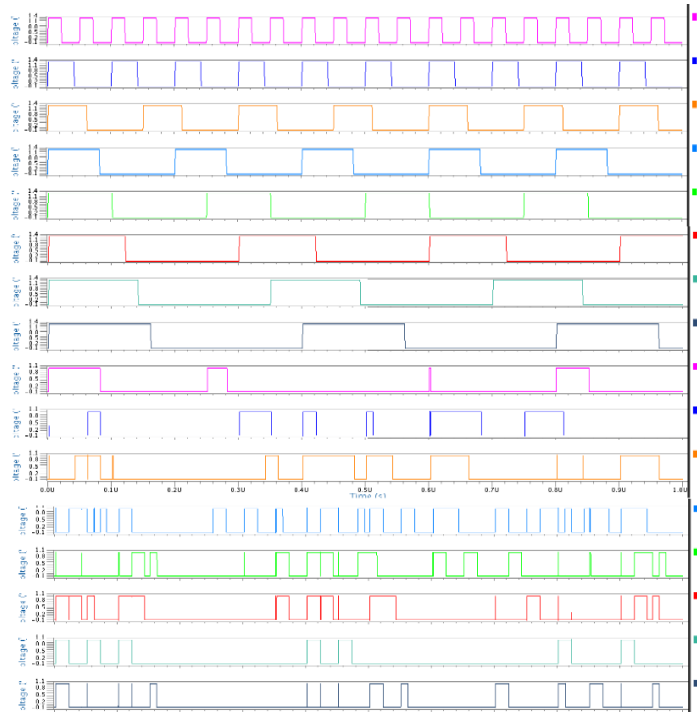


Fig -4: Output waveforms of conventional Braun multiplier

Table I. shows the list of time period and pulse width given to the inputs of the multiplier in test bench.

Table -1: Inputs for Conventional Braun Multiplier

Inputs	Pulse Width(ns)	Time period(ns)
X0	50	20
X1	100	40
X2	150	60
X3	200	80
Y0	250	100
Y1	300	120
Y2	350	140
Y3	400	160

4. KOGGE-STONE ADDER

Kogge-stone adder(KSA) is a parallel prefix adder. Carry calculation is as fast as carry generation time $O(\log n)$. Braun multiplier is designed to reduce area and increase its speed with Kogge-Stone adder compared to conventional Braun multiplier, which is used in most modern DSPs. In [6], 4-bit Braun Multiplier with KSA is implemented on FPGA using HDL Verilog which reduces delay. In the last stage, where its functionality is similar to RCA, the delay of the whole multiplier depends mostly on the delay of Full adder array. The RCA-produced Delay can be reduced by using one of the "KOGGE STONE ADDER" Parallel Prefix Adders[4].

A. Illustration of KSA using an example of 3 bit

A=011 and B = 100 are inputs

a) I Step: Pre processing

Computation of generate and propagate signals $P_i = A_i \text{ XOR } B_i$ and $G_i = A_i \text{ AND } B_i$

- $P_3 = 0 \text{ XOR } 1 = 1;$
- $P_2 = 1 \text{ XOR } 0 = 1;$
- $P_1 = 1 \text{ XOR } 0 = 1;$
- $G_3 = 0 \text{ AND } 1 = 0;$
- $G_2 = 1 \text{ AND } 0 = 0;$
- $G_1 = 1 \text{ AND } 0 = 0;$

b) II Step: Carry look ahead network

Computation of carries corresponding to each bit $P_{i:j} = P_{i:k+1} \text{ AND } P_{k:j}$ and $G_{i:j} = G_{i:k+1} \text{ OR } (P_{i:k+1} \text{ AND } G_{k:j})$

$P_{3:2} = P_3 \text{ AND } P_2; 1 \text{ AND } 1 = 1;$
 $G_{3:2} = G_3 \text{ OR } (P_3 \text{ AND } G_2); 0 \text{ OR } (1 \text{ AND } 0) = 0;$

$P_{2:1} = P_2 \text{ AND } P_1; 1 \text{ AND } 1 = 1;$

$G_{2:1} = G_2 \text{ OR } (P_2 \text{ AND } G_1); 0 \text{ OR } (1 \text{ AND } 0) = 0;$

$P_{3:1} = P_{3:2} \text{ AND } P_1; 1 \text{ AND } 1 = 1;$

$G_{3:1} = G_{3:2} \text{ OR } (P_{3:2} \text{ AND } G_1); 0 \text{ OR } (1 \text{ AND } 0) = 0;$

c) III Step: Post processing

Computation of sum

bits $S_i = P_i \text{ XOR } C_{i-1}$

$S_1 = P_1 \text{ XOR } 0; 1 \text{ XOR } 0 = 1;$

$S_2 = P_2 \text{ XOR } G_1; 1 \text{ XOR } 0 = 1;$

$S_3 = P_3 \text{ XOR } G_2; 1 \text{ XOR } 0 = 1;$

$C_{out} = G_{3:1}; 0;$

Sum Bits : $C_{out} S_3 S_2 S_1 : 0111;$

B. 3-Bit Kogge Stone Adder with 12T

Fig. 5 is a 3-bit CMOS Kogge-stone adder. It is made of 6 XOR gates (6*12), 10 AND gates and 5 OR gates with total of 510 transistors.

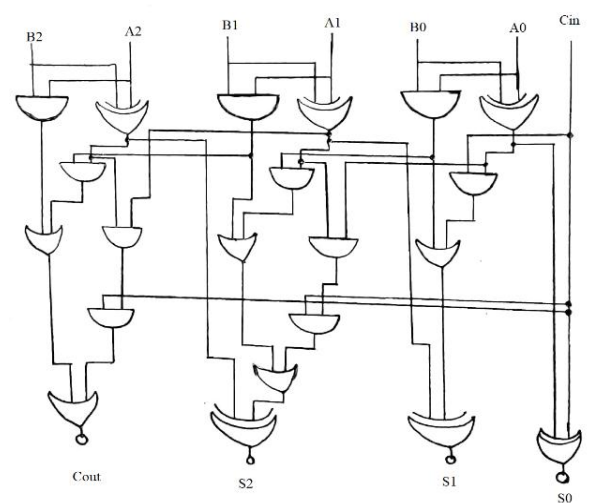


Fig -5: Schematic of 3-bit KSA

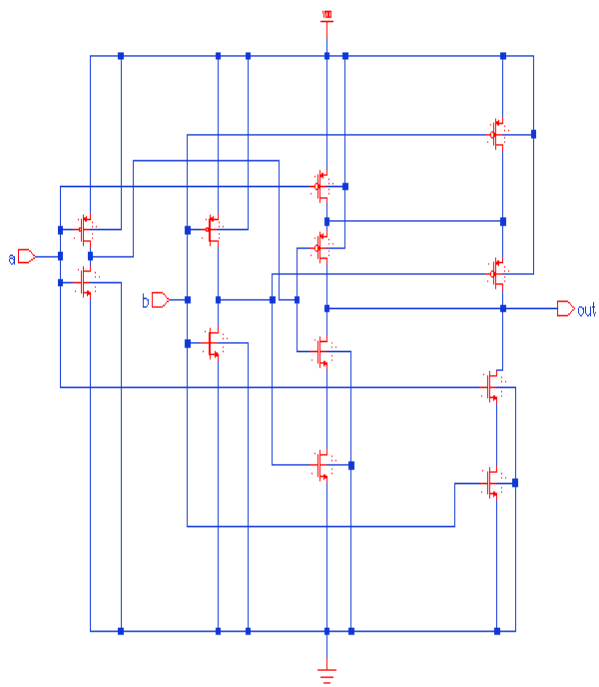


Fig -6: Schematic of 12T XOR

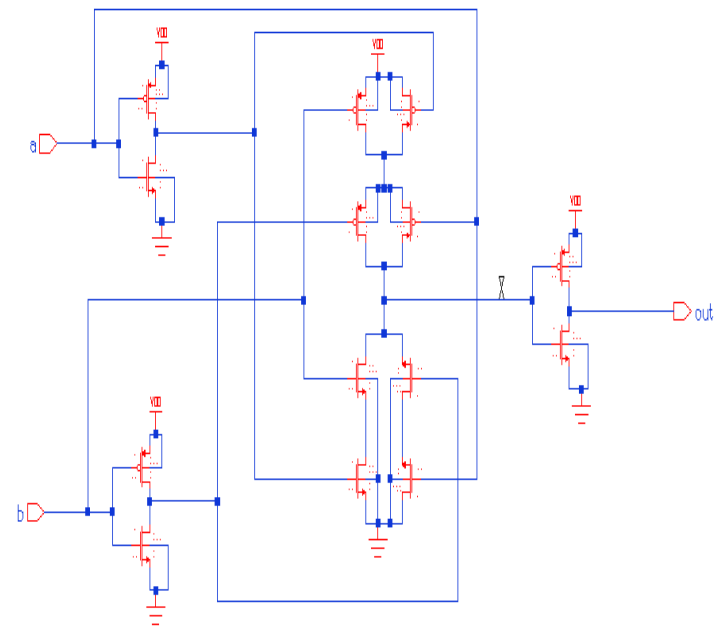


Fig -8: 14T XOR used in KSA

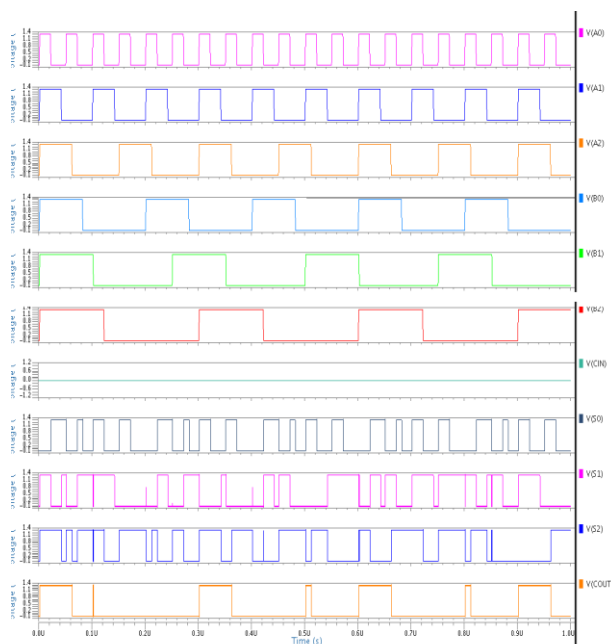


Fig -7: Waveforms of 3-bit KSA

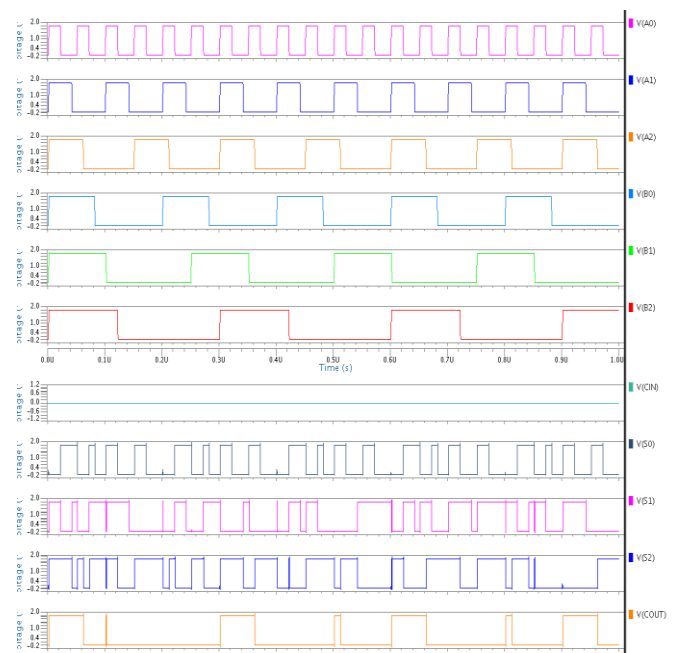


Fig -9: Waveforms of KSA with 14T XOR gate

C. Kogge-Stone Adder with 14T XOR gates

Fig. 5 shows the 3-bit KSA used. Fig. 8 is 14T XOR gate design and KSA output results are shown in Fig. 9. 14T KSA requires 6 XOR gates (6*14), 10 AND gates and 5 OR gates.

D. Kogge-Stone Adder with 22T XOR gates

Fig. 10 shows the schematic of 22T XOR gate used in KSA. Fig. 11 shows the output waveforms of the 22T XOR gate KSA.

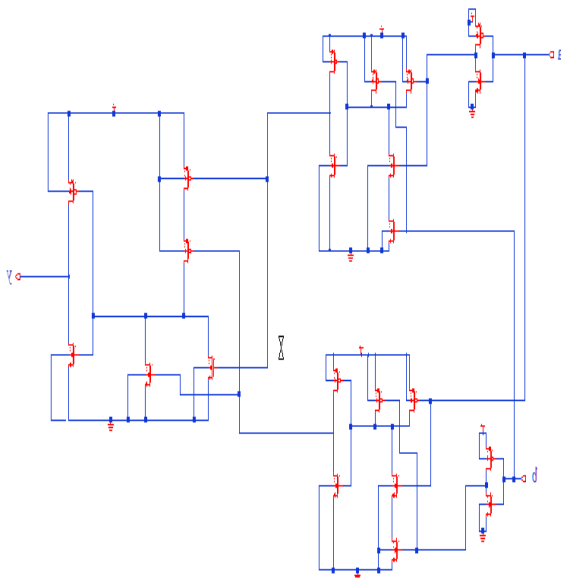


Fig -10: Schematic of 22T XOR gate

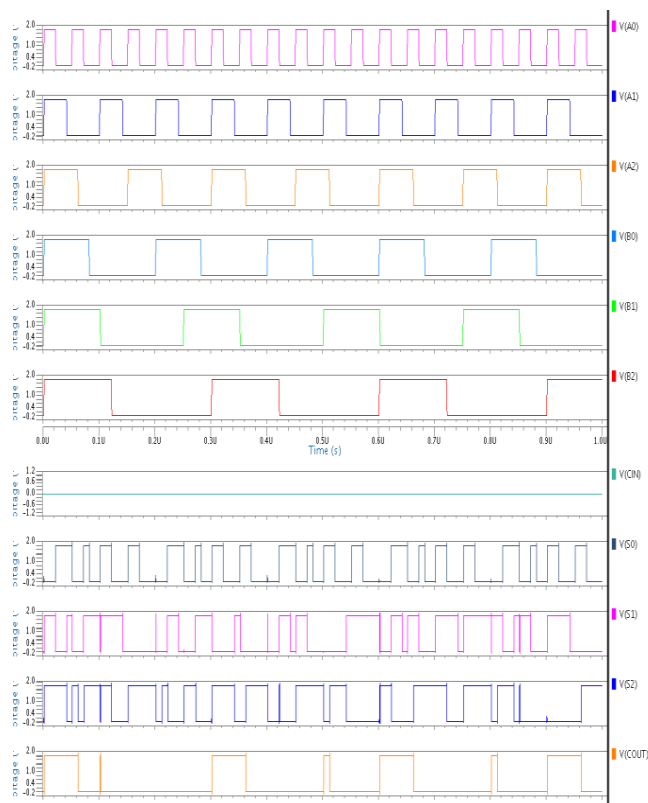


Fig -11: Waveforms of KSA with 22T XOR gate

E. Comparison of Results of KSA

For KSA with 14 T XOR design delay, the power requirement and the transistor count is lower compared to the 22 T XOR KSA design. Table II shows a comparison between 12T, 14T, and 22T KSA.

Table -2: Comparison Results of KSA

KSA	Delay (180nm)	Delay (130nm)	Transistor count
22T-XOR	311.50ps	101.09ps	222
14T-XOR	21.285ns	152.91ps	174
12-XOR	21.363ns	21.03ns	162

5. SIMULATION RESULTS OF BRAUN MULTIPLIER WITH KOGGE-STONE ADDER

The proposed multiplier's block diagram makes use of a three bit KSA in 4th stage of Braun multiplier. This proposed multiplier is implemented using Mentor graphics 180nm technology. To design 4 bit (n bit) Braun, it needs 3bit (n-1 bit) KSA. KSA is designed with 22T XOR, 14T XOR and 12T XOR gates. The circuit shown in the Fig. 12 is modified Braun multiplier with 4x4 input and 8 output pins. of 8 outputs pins, P₄-P₇ pins cause more delay in the output. Because delay is mainly due to these bits, this stage is replaced with KSA. Not only P₄-P₇ pins stage but also other stages can be designed using KSA. But the area increases and system defects create problem for a couple of stages.

The circuit shown in the Fig. 13 is test bench for proposed Braun multiplier made using CMOS technology. The output of Braun multiplier with KSA is represented in Fig. 13. Table III represents the comparison between conventional Braun multiplier and the new design.

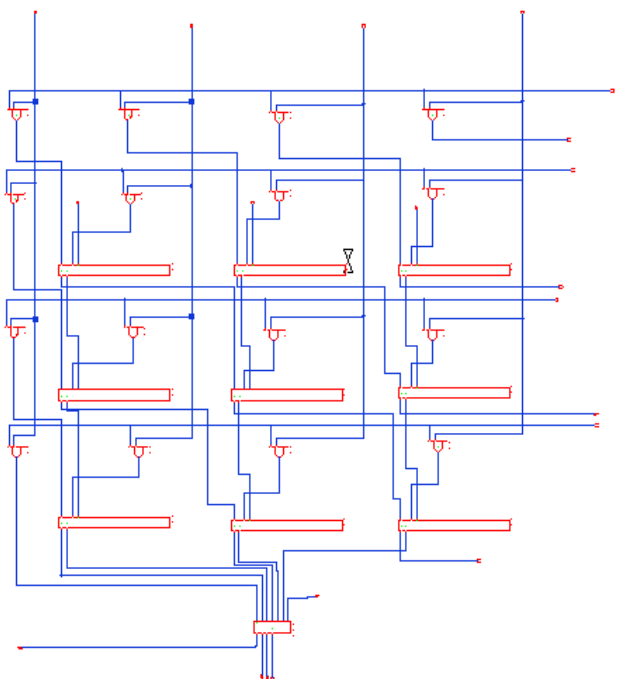


Fig -12: Schematic of proposed Braun multiplier

Table -3: Comparison of Conventional and Proposed designs of Braun Multiplier

4-bit Braun multiplier	Delay (180nm)	Delay (130nm)
Conventional Design	535.72ps	221.32ps
Proposed design with KSA of 22T XOR	134.72ps	124.48ps
Proposed design with KSA of 14T XOR	226.63ps	124.48ps
Proposed design with KSA of 12T XOR	236.71ps	124.48ps

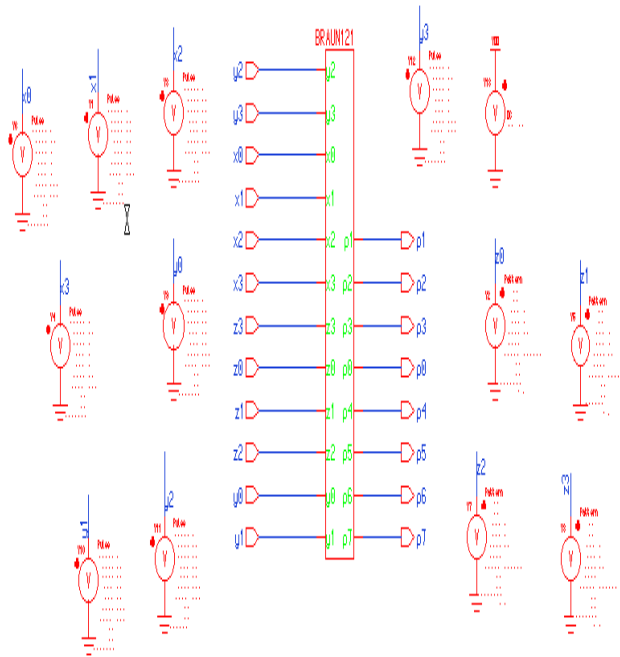


Fig -13: Test bench of proposed Braun multiplier

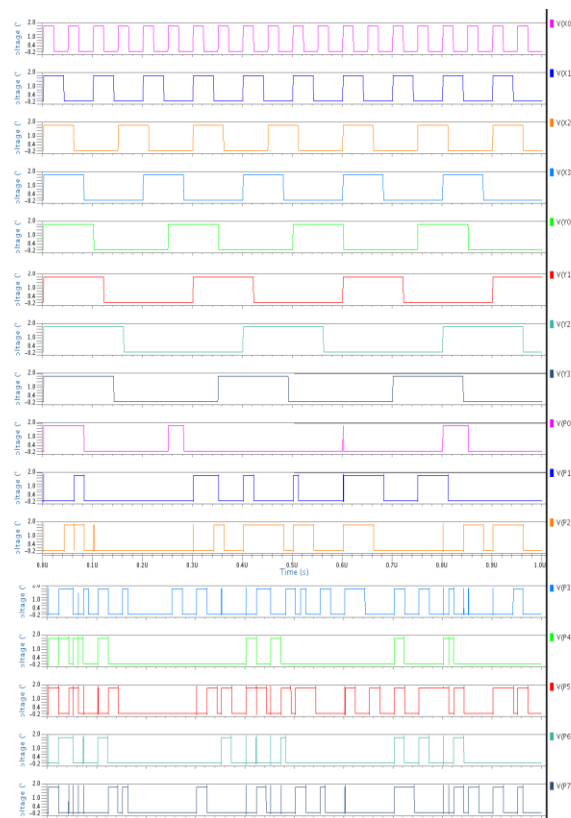


Fig -14: Output of Proposed Braun multiplier with KSA of 12T XOR

6. CONCLUSION

A conventional Braun multiplier and Braun multiplier with new design is designed and compared in terms of delay. Conventional Braun multiplier is designed with Full adder which consists of 28 transistors and AND gates comprising of 6 Transistors. The new design consists of 16 AND gates, 9 Full adders, one Kogge-stone adder (KSA). The KSA is designed with 3 models those are with 22T, 14T, 12T XOR gates. The delay reduced from Conventional Braun multiplier to new design of Braun multiplier with 22T KSA is 401ns, 309.09ns for Braun multiplier with 14T KSA and 299.01ns for Braun multiplier with 12T for 180nm and 96.84ns for 130nm technology.

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