

# Active Directory Golden Ticket Attack Detection

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**Abstract** – Active Directory is one of the most popular identity and access management systems for corporate networks, which makes it an extremely important component of companies' IT infrastructure and, at the same time, an attractive target for hackers. Ticket forging attacks, such as Golden and Silver Ticket are considered to be especially dangerous, as they allow malicious actors to escalate their privileges and gain an unauthorized access to services and resources on the network. Therefore, quick and accurate detection of forged tickets is critical for network protection and timely response on potential threats. This article discusses a method of detecting the malicious Golden Ticket activity based on the traffic and attack signature analysis in an Active Directory-based network.

**Key Words:** Kerberos, Active Directory, forged tickets, domain controller, authorization

## 1. INTRODUCTION

Active Directory (AD) is Microsoft's directory service used for managing resources and access on the network. It serves as a centralized database that stores information about users and their permissions, groups, computers, and other network resources. Active Directory Domain Services (AD DS) stands as the fundamental element within Active Directory, facilitating user authentication and resource access across the network. Active Directory arranges objects in a hierarchical structure, enabling different Domain Services to interact with these objects and granting users the ability to access or oversee them [1].

Active Directory was developed as a solution to the challenges of centralized management and resource protection, which are especially critical in large organizations. Prior to the advent of AD, controlling a multitude of workstations, services, user accounts, and permissions, as well as managing service servers, was a complex and time-consuming task. The absence of a centralized management system resulted in repetitive actions and unreasonably high efforts.

With AD, administrators gained the ability to manage the entire system from a single entry point, significantly accelerating the management process and saving a considerable amount of time. Today, AD is integrated into the infrastructure of more than 90% of companies listed in the Global Fortune 1000 [2].

Despite its widespread use, Active Directory networks often remain vulnerable. As the need to adapt to modern remote work conditions grows, security specialists often have to choose between smooth user experience and higher protection, which often means more constraints for remote workers comparing to those, who traditionally works from office. The extensive use of AD, coupled with inadequate protection, makes it a prime target for cybercriminals, Advanced Persistent Threat (APT) groups, and threat actors seeking unauthorized access to organizations' resources and company data. The whitepaper [3] provides an overview of key APT features, including APT terminology, lifecycle, techniques, types of targets, comparison with malware, detection and protection measures. General Active Directory security concerns are discussed in [4] and [5], while an analysis of cyberattack methods is presented in [5]. Techniques and tools commonly employed by hackers are discussed in [6] and [7].

One of the most perspective Active Directory attack vector targets its authentication protocol, Kerberos [8]. Kerberos authentication relies on time-restricted cryptographic messages known as tickets, which verify the user's identity to a target server without transmitting credentials over the network or storing them locally. While this mechanism offers numerous advantages for both end users and system administrators, it is not without its security flaws and is vulnerable to forged ticket attacks: Golden Ticket [9] and Silver Ticket [10]. These attacks enable an intruder to establish a domain persistence on the post-exploitation phase, move laterally within the network or escalate their privileges, granting access to restricted and confidential resources while bypassing standard authentication methods. Consequently, Active Directory networks require continuous monitoring and analysis to detect changes in the environment promptly. This vigilance is essential to reduce potential risks and losses.

This article provides an overview of the principles behind Golden and Silver Ticket attacks, along with network signatures and algorithm concept that can be employed to identify and thwart these attacks, thereby preventing potential harm.

## 2. KERBEROS AUTHENTICATION ALGORITHM

The main components of the Kerberos authentication system are [11]:

- Authentication Service (AS): This service is responsible for authenticating clients on the network.
- Ticket-Granting Service (TGS): The TGS is responsible for issuing Service Tickets (ST) and Ticket-Granting Tickets (TGT) to end-users. These tickets are essential for end-users to access any Kerberos-enabled services in the Active Directory environment.
- Key Distribution Center (KDC): The KDC runs from the krbtgt account on the Domain Controller and encapsulates both AS and TGS.

The Kerberos authentication process consists of the following steps:

### 1. AS Exchange

#### a. KRB\_AS\_REQ

In this step, the user initiates an authentication process by providing their username and password during login. After the user submits their input, the KRB\_AS\_REQ request is sent to the AS. The request contains user's information (SID, name, group membership, etc.), along with a timestamp encrypted with a key derived from their credentials. The AS checks the existence of the supplied username in its database by querying the NTDS.dit file, which stores all the Active Directory data. If the username is found, AS uses the corresponding password hash to attempt to decrypt the provided timestamp. If the attempt is successful, the KDC is ensured that the user is who they claim to be. It then generates a unique session key for communication with the Ticket-Granting Service. This session key is tied to the user and restricted in time.

#### b. KRB\_AS\_REP

The KDC's response includes the session key for TGS, encrypted with the user's password hash, and the TGT ticket. The TGT contains various information, including the username, validity period, a copy of the generated session key, and a

Privilege Attribute Certificate (PAC). The TGT itself is encrypted with the krbtgt key, ensuring that only the KDC can decrypt it and access its contents.

With the TGT in their possession, the user is now considered authenticated within the environment and can communicate with the TGS.

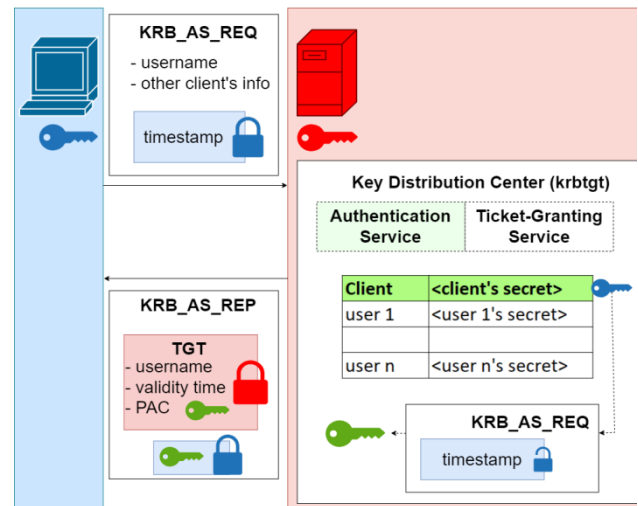


Figure 1: Kerberos AS Exchange

### 2. TGS Exchange

#### a. KRB\_TGS\_REQ

If an authenticated user wants to use a specific service within the environment, they send a request to the TGS. This request includes the Ticket-Granting Ticket (TGT), obtained during KRB\_AS\_REQ, the Service Principal Name (SPN) of the target service, and an authenticator containing the username and timestamp. The authenticator is encrypted with the TGS session key, which was obtained during the AS Exchange.

The TGS then compares the data within the TGT and the authenticator. The KDC extracts the contents of the TGT, including the user's name and the corresponding session key. Using this extracted session key, the KDC attempts to decrypt the authenticator. If the decryption is successful and the data within the authenticator matches the data within the TGT, the user's identity is confirmed. The KDC is assured that the requester

possesses the TGT and knows the specific session key.

b. KRB\_TGS\_REQ

The KDC responds by providing the user with the necessary information to request access to the service. This message includes a Service Ticket (ST) and a new session key. This new session key is time-restricted and valid only for communication between the user and the target server, whose Service Principal Name (SPN) was specified during the KRB\_TGS\_REQ.

The ST contains a copy of the session key, the SPN of the target service, the username, and the Privilege Attribute Certificate (PAC). To secure this ticket, it is encrypted with the target service's key, ensuring that only the target service and the KDC can decrypt, read, and modify its contents. Additionally, both the ticket and the session key are encrypted again, this time using the session key that facilitated communication between the client and the TGS.

Upon receiving the response, the client can decrypt only the first layer, which is encrypted with the TGS session key. This allows the client to extract the new session key for communication with the target server and the encrypted ST.

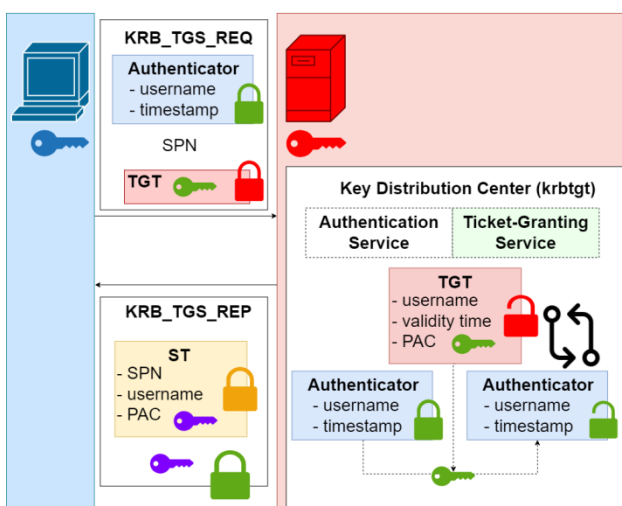


Figure 2: Kerberos TGS Exchange

3. Client-Server Exchange

a. KRB\_AP\_REQ

The client generates a new authenticator, encrypts it with the new session key received in the KRB\_TGS\_REQ, and sends it to the service server along with the ST. The service server receives the ST and can decrypt it using its own secret. The session key included within the ST is then used to decrypt the authenticator.

By comparing the data within the ST and the authenticator, the server can verify the client's authenticity. If the data within the ST matches the data in the authenticator, the client is authorized to use the service.

b. KRB\_AP\_REP

In certain cases, if the client has set the mutual authentication flag in the KRB\_AP\_REQ, the server will use the session key to encrypt the timestamp sent by the client. This encrypted timestamp is then sent back to the client. This message is known as the KRB\_AP\_REP. The client decrypts the timestamp and compares it with the original timestamp they sent in the KRB\_AP\_REQ. If these timestamps match, the client is assured of the server's authenticity, and communication can proceed.

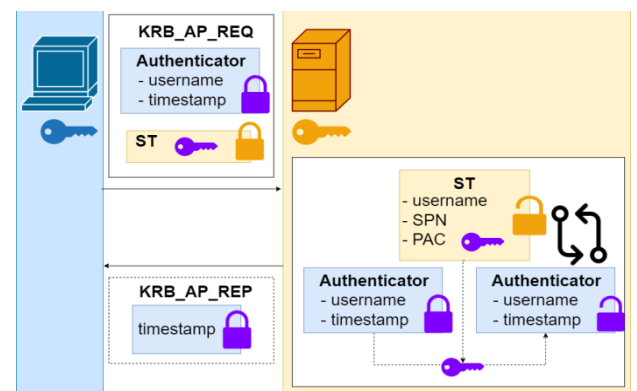


Figure 3: Kerberos Client-Server Exchange

3. GOLDEN TICKET ATTACK

The Golden Ticket attack is aimed at falsifying the Ticket Granting Ticket (TGT), which is crucial for domain authentication. Notably, Kerberos itself lacks inherent authorization mechanisms [12]. In an Active Directory (AD) environment, authorization is accomplished through

the Privilege Attribute Certificate (PAC) stored within the TGT and Service Tickets (ST). PAC within TGT is used for STs creation. When authorizing a client, the server running the service examines various factors, including group membership, user roles, and rights mentioned in the PAC.

In the process of forging the TGT, an attacker can specify arbitrary PAC contents, which serve as the basis for all subsequent service tickets. The scope of this attack encompasses the entire domain infrastructure. With a forged TGT, an attacker gains the ability to generate service tickets for any domain services with the desired rights, effectively compromising the entire domain. The service tickets created based on the manipulated TGT can have arbitrary validity times, not constrained by domain policies, as the Key Distribution Center (KDC) inherently trusts the TGT. This trust exists because, in the typical Kerberos authentication process, the KDC issues TGTs itself without questioning their authenticity.

According to official Microsoft documentation [13], PAC validation is typically not applied to TGT tickets, unless they are older than 20 minutes. TGTs are encrypted and signed with the krbtgt account hash. Therefore, to forge such a ticket, an attacker must compromise the krbtgt hash.

The typical attack flow is as follows:

1. An attacker successfully compromises the krbtgt account hash.
2. The attacker creates their own TGT, similar to one that can be found in KRB\_AS\_REP. Consequently, they can set their own arbitrary privileges within the PAC.
3. The ticket created by the attacker is signed with the compromised service hash. The attacker creates a KRB\_TGS\_REQ from scratch and obtain an ST for any domain service.
4. The falsified KRB\_TGS\_REQ is sent to the DC and based on the PAC within provided TGT the DC issues an ST. If the PAC indicates that the user is a domain administrator, they will be granted administrative privileges

Subsequent communication with application servers follows the normal workflow: the service successfully decrypts the ST, extracts the session key, decrypts the authenticator, and grants users privileges according to their PAC. With a forged TGT, an attacker possesses the desired privileges on all domain servers. An additional threat posed by this attack is that, by default, the krbtgt password is almost never changed, making this attack especially effective for achieving long-term domain persistence.

## 4. GOLDEN TICKET ATTACK DETECTION

### 4.1 Background

Because ticket forging attacks are typically employed during the post-exploitation phase to establish persistence [3], many companies have designed their security systems to thwart these attacks by implementing robust security procedures and policies, along with deploying controls and protections at earlier stages. However, if an attack occurs, an attacker can remain in the network undetected for a long time, making these attacks extremely effective in gaining a foothold in the domain, and there is currently no universal or complete solution to protect against these attacks. That's why it is so important to timely detect the occurred breach. Some of the common approaches when detecting Golden Ticket attacks are:

- Monitoring for anomalies in Windows logs (Event IDs 4624, 4672 or 4634) [14].
- Analyzing Windows security logs for suspicious activity [15]
- Monitoring Kerberos tickets lifetime [14]
- Monitoring of the outdated encryption algorithms used for Kerberos tickets [16]

Although there are ways to detect ticket forging attacks, many of them can be circumvented by manipulating requests, or existing detection methods can respond to legitimate activity and produce a large number of false positives. The offered method relies on the sequence of network message exchanges and is described in-detail below.

### 4.2 Theoretical basis

The most effective way to detect Golden Tickets is by correlating TGS requests with AS requests. In legitimate activities, AS requests always precede TGS requests. The communication between the client and the server is illustrated in Figure 4.

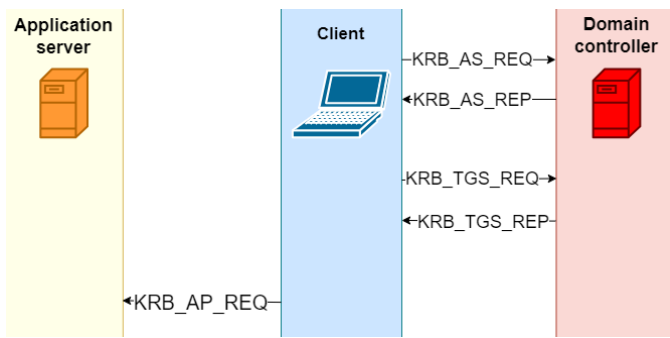


Figure 4: Legitimate Kerberos messages sequence

In the Golden Ticket attack scenario, the KRB\_AS\_REQ and KRB\_AS\_REP messages are absent because the client does not request a TGT from the domain controller but creates it from scratch. Client generates KRB\_TGS\_REQ in order to obtain the service ticket and network communication looks as shown on the Figure 5.

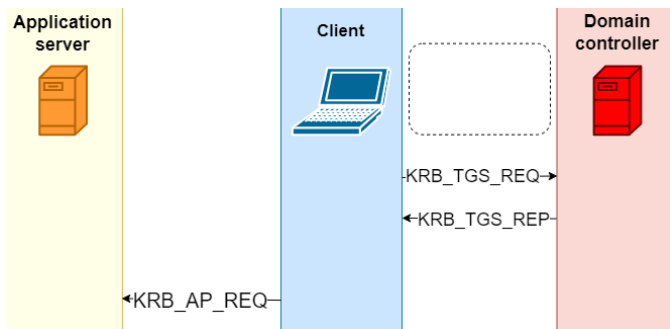


Figure 5: Malicious Kerberos messages sequence

Consequently, the first Golden Ticket attack signature is a violation of the normal message exchange sequence. However, this signature alone is insufficient, as it does not cover the following scenarios:

1. Scenario 1
  - a. An attacker authenticates in the network, performs KRB\_AS\_REQ, and receives TGT
  - b. The attacker either forges another TGT or decrypts the current one and alters the data inside it
2. Scenario 2
  - a. A legitimate user authenticates on the network and obtains a TGT.
  - b. An attacker tampers TGT and sets the same username as the username of the legitimate user.

Both scenarios are keeping the normal Kerberos message sequence unaltered, so the initial signature will not be

triggered. As mentioned earlier, TGTs are transmitted on the network in an encrypted form and appear as text lines. Encrypted TGTs contain a timestamp and a session key, making each issued TGT unique. Therefore, if any data within the TGT is altered, it impacts the entire text line transmitted on the network. Hence, the second attack signature involves a discrepancy between the TGT issued in KRB\_AS\_REP and the one presented in KRB\_TGS\_REQ.

The attack detection algorithm involves network traffic analysis and logging of all TGT tickets issued by the domain controller. It then compares the TGTs presented by users with the entries in the table. If the presented TGT is absent from the table, it indicates that the KDC had never issued it, and it was created by an attacker from scratch.

### 4.3 Practical implementation

For the attack simulation, we created an Active Directory virtual network with the following components:

- Domain Controller: Windows Server 2019
- Domain name: SAMPLE.local
- Client: Windows 10 Enterprise
- Users:
  - Administrator@SAMPLE.local – SAMPLE.local domain administrator
  - jdoe@SAMPLE.local – non-privileged user of the SAMPLE.local domain

The Domain Controller hosts an SMB server with resources accessible only to domain administrators. This SMB server can be accessed via the path \\MLNK-DC\hackme and contains a test file called congrats.txt. Non-privileged users do not have access to this shared SMB resource. In the attack scenario, this SMB file represents the protected resource that attackers are attempting to access.

Figure 6 illustrates the attack environment.

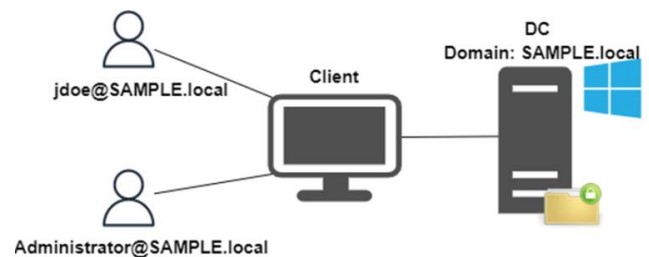


Figure 6: Virtual network for Golden Ticket simulation

The virtual network was created using VMware as a hypervisor. Each workstation is a separate virtual machine

with a NAT adapter. Virtual machines with NAT function similarly to real computers when connecting to the internet via a router, which, in this case, is the network kernel of VMware.

For carrying out the attack, we utilized a tool called Mimikatz. Mimikatz is an open-source application that allows users to view and save authentication data, including Kerberos tickets. It's often used by malicious actors for credential theft and privilege escalation. As a result, it's frequently detected and removed by endpoint protection systems and antivirus software.

It's essential to note that ticket forging attacks, like the Golden Ticket attack, are typically used during the post-exploitation phase to establish domain persistence. This implies that for a Golden Ticket attack to succeed, an attacker must first gain access to the network and compromise the krbtgt password hash. There are various methods to compromise the krbtgt hash, with one common approach being to access the domain controller's database, found in the C:\Windows\NTDS\ntds.dit file. In our case, we used the DCsync attack, which utilizes the Directory Replication Service (DRS), to obtain passwords from the NTDS.DIT file.

Normally, to execute the DCsync attack, an attacker would require rights such as 'Replicating Directory Changes' or 'Replicating Directory Changes All.' These rights are typically granted to users in the Administrators, Domain Admins, and Enterprise Admins groups. However, they may also be present due to system administrator negligence or incorrectly assigned privileges. The specific methods for compromising the krbtgt password hash are beyond the scope of this paper.

The network traffic was recorded with the Wireshark tool, which was running on the DC.

To facilitate comparison and analysis of the network traffic, the following scenarios were simulated:

- Scenario 1 (Legitimate activity)

Administrator logs in their account and attempts to access the protected SMB share as the Figure 7 demonstrates. Command to list the share directory is as follows:

```
dir \\MLNK-DC\hackme
```

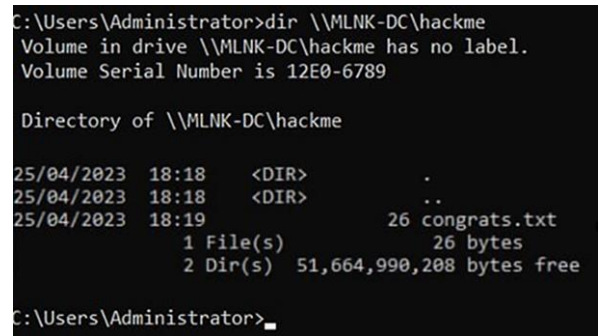


Figure 7: Administrator accesses SMB share

Administrator is a privileged user so they are allowed to access the resource.

- Scenario 2 (Golden Ticket activity)

Jdoe logs into their account and initially attempts to access the restricted resource using the TGT issued by the domain controller

Command to list available Kerberos tickets with Mimikatz is as follows:

```
kerberos::list
```

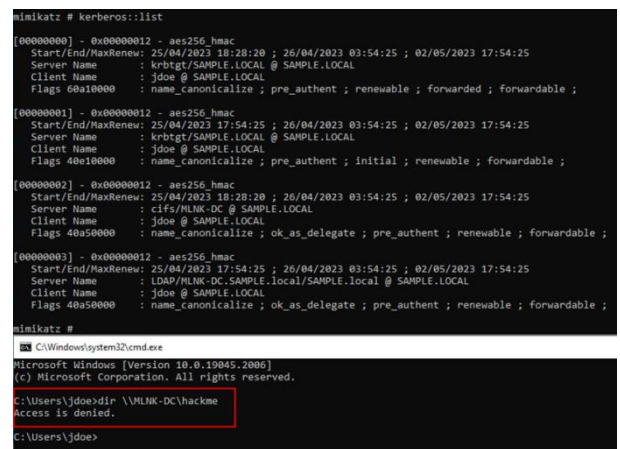


Figure 8: Jdoe is not privileged enough to access SMB share

Since Jdoe is a non-privileged domain user without permissions to access the protected SMB share using the TGT issued by the domain controller, they proceed to conduct the Golden Ticket attack. Jdoe utilizes the falsified TGT to gain new privileges for the current session, as illustrated in Figure 9.

Command to generate golden ticket and then inject it in the current session are as follows:

```
kerberos::golden /user:attacker
/domain:SAMPLE.local /sid:S-1-5-21-2593423864-4051040445-2235811965
/krbtgt:65917334c795c32fd7656534a9bc9ab2
/ticket:gld.tck /ptt
```

```
mimikatz 2.2.0 x64 (x64)
Flags 40a50000 : name_canonicalize ; ok_as_delegate ; pre_authent ; renewable ; forwardable ;
mimikatz # kerberos::golden /user:attacker /domain:SAMPLE.local /sid:S-1-5-21-2593423864-4051040445-2235811965 /krbtgt:65917334c795c32fd7656534a9bc9ab2 /ticket:gld.tck /ptt
User : attacker
Domain : SAMPLE.local (SAMPLE)
SID : S-1-5-21-2593423864-4051040445-2235811965
User Id : 500
Groups Id : *513 512 520 518 519
ServiceKey : 65917334c795c32fd7656534a9bc9ab2 - rc4_hmac_nt
Lifetime : 25/04/2023 18:30:54 ; 22/04/2023 18:30:54 ; 22/04/2023 18:30:54
-> Ticket : ** Pass The Ticket **
* PAC generated
* PAC signed
* EncTicketPart generated
* EncTicketPart encrypted
* KrbCred generated
Golden ticket for 'attacker@SAMPLE.local' successfully submitted for current session
mimikatz # kerberos::list
[00000000] - 0x00000017 - rc4_hmac_nt
Start/End/MaxRenew: 25/04/2023 18:30:54 ; 22/04/2023 18:30:54 ; 22/04/2023 18:30:54
Server Name : krbtgt/SAMPLE.local @ SAMPLE.local
Client Name : attacker @ SAMPLE.local
Flags 40e00000 : pre_authent ; initial ; renewable ; forwardable ;
Flags 40e00000 : pre_authent ; initial ; renewable ; forwardable ;
mimikatz #
```

Figure 9: Jdoe forces TGT and uses it for current session

With the forged TGT the non-privileged user Jdoe is able to access the protected resource, as demonstrated on the Figure 10.

```
mimikatz # kerberos::tgt
kerberos: TGT of current session :
Start/End/MaxRenew: 25/04/2023 18:30:54 ; 22/04/2023 18:30:54 ; 22/04/2023 18:30:54
Service Name (00) : krbtgt : SAMPLE.local ; @ SAMPLE.local
Target Name (-) : @ SAMPLE.local
Client Name (01) : attacker ; @ SAMPLE.local
Flags 40e00000 : pre_authent ; initial ; renewable ; forwardable ;
Session key : 0x00000017 - rc4_hmac_nt
Ticket : 0x00000017 - rc4_hmac_nt ; kvno = 0 [...]
** Session key is NULL! It means allowtgsessionkey is not set to 1 **
mimikatz #
C:\Windows\system32\cmd.exe
Microsoft Windows [Version 10.0.19045.2006]
(c) Microsoft Corporation. All rights reserved.
C:\Users\jdoe>dir \\VULNK-DC\hackme
Access is denied.
C:\Users\jdoe>dir \\VULNK-DC\hackme
Volume in drive V:\VULNK-DC\hackme has no label.
Volume Serial Number is: 2108-6790
Directory of V:\VULNK-DC\hackme
25/04/2023 18:18 <DIR> .
25/04/2023 18:18 <DIR> ..
25/04/2023 18:19 20 congrats.txt
25/04/2023 18:19 1 File(s) 20 bytes
25/04/2023 18:19 2 Dir(s) 51,664,998,288 bytes free
C:\Users\jdoe>
```

Figure 10: Jdoe successfully accesses SMB share bypassing authorization controls

### 4.3 Results analysis

Figures 11 and 12 depict intercepted network traffic within Scenario 1. This activity is legitimate, none of the discussed above signatures were triggered. The message sequence is consistent, and the encrypted ticket provided in TGS\_REQ matches the ticket issued by the domain controller in AS\_REP.

```
-- 0... 192.168.226.141 192.168.226.144 KRB5 359 AS-REQ
-- 0... 192.168.226.144 192.168.226.141 KRB5 1635 AS-REP
-- 0... 192.168.226.141 192.168.226.144 KRB5 79 TGS-REQ
-- 0... 192.168.226.144 192.168.226.141 KRB5 1605 TGS-REP
-- 0... 192.168.226.141 192.168.226.144 KRB5 247 TGS-REQ
> Frame 13: 1635 bytes on wire (13080 bits), 1635 bytes captured (13080 bits) on interface NDevice\WIFI
> Ethernet II, Src: VMware_e3:ad:7f (00:0c:29:e3:ad:7f), Dst: VMware_cd:bb:41 (00:0c:29:cd:bb:41)
> Internet Protocol Version 4, Src: 192.168.226.144, Dst: 192.168.226.141
> Transmission Control Protocol, Src Port: 88, Dst Port: 51840, Seq: 1, Ack: 306, Len: 1581
Kerberos
Record Mark: 1577 bytes
as-rep
pvno: 5
msg-type: krb-as-rep (11)
padata: 1 item
crealm: SAMPLE.LOCAL
name
ticket
tko-vno: 5
realm: SAMPLE.LOCAL
sname
enc-part
etype: eTYPE-AES256-CTS-HMAC-SHA1-96 (18)
kvno: 2
cipher: 8ab0aba64322743d897c649e3a736abb7e6fa61a5dff7be5bcc0779e1d0051607d08e97d...
```

Figure 11: TGT issued to the user in AS\_REP

```
-- 0... 192.168.226.141 192.168.226.144 KRB5 359 AS-REQ
-- 0... 192.168.226.144 192.168.226.141 KRB5 1635 AS-REP
-- 0... 192.168.226.141 192.168.226.144 KRB5 79 TGS-REQ
-- 0... 192.168.226.144 192.168.226.141 KRB5 1605 TGS-REP
-- 0... 192.168.226.141 192.168.226.144 KRB5 247 TGS-REQ
-- 0... 192.168.226.144 192.168.226.141 KRB5 1717 TGS-REP
-- 0... 192.168.226.141 192.168.226.144 DCERPC 610 Bind: call_id: 2, Fragment: Single, 3 context items:
-- 0... 192.168.226.141 192.168.226.144 DCERPC 339 Bind_ack: call_id: 2, Fragment: Single, max_xmit: 584
-- 0... 192.168.226.141 192.168.226.144 DCERPC 274 Alter_context: call_id: 2, Fragment: Single, 1 context
-- 0... 192.168.226.141 192.168.226.144 LDAP 385 bindRequest(75) "sasl"
-- 0... 192.168.226.144 192.168.226.141 LDAP 265 bindResponse(75) success
-- 0... 192.168.226.141 192.168.226.144 LDAP 501 bindRequest(79) "sasl"
-- 0... 192.168.226.144 192.168.226.141 LDAP 265 bindResponse(79) success
-- 48... 192.168.226.141 192.168.226.144 KRB5 219 TGS-REQ
-- 48... 192.168.226.144 192.168.226.141 KRB5 1678 TGS-REP
-- 48... 192.168.226.141 192.168.226.144 KRB5 1494 TGS-REQ
-- 48... 192.168.226.144 192.168.226.141 KRB5 1522 TGS-REP
-- 48... 192.168.226.141 192.168.226.144 SMB2 422 Session Setup Request
-- 48... 192.168.226.144 192.168.226.141 SMB2 315 Session Setup Response
enc-part
etype: eTYPE-AES256-CTS-HMAC-SHA1-96 (18)
kvno: 2
cipher: 8ab0aba64322743d897c649e3a736abb7e6fa61a5dff7be5bcc0779e1d0051607d08e97d...
authenticator
PA-DATA pA-PAC-OPTIONS
req-body
padding: 0
kdc-options: 40810000
realm: SAMPLE.LOCAL
sname
name-type: kRB5-NT-SRV-INST (2)
sname-string: 2 items
SNameString: cifs
SNameString: MLNK-DC
```

Figure 12: TGT presented by the user in TGS\_REQ

Figures 13 and 14 showcase intercepted network traffic within Scenario 2, which is an attack scenario. Initially, the network activity may appear legitimate, given the consistent message sequence. However, upon closer examination of the packets, it becomes evident that the ticket issued by the domain controller in AS\_REP does not match the ticket submitted by the user in TGS\_REQ. This discrepancy serves as an attack signature.

```

- 0.. 192.168.226.141 192.168.226.144 KRBS 350 AS-REQ
- 0.. 192.168.226.144 192.168.226.141 KRBS 1541 AS-REP
- 0.. 192.168.226.141 192.168.226.144 KRBS 1456 TGS-REQ
- 0.. 192.168.226.144 192.168.226.141 KRBS 1483 TGS-REP
- 0.. 192.168.226.141 192.168.226.144 KRBS 164 TGS-REQ
- 0.. 192.168.226.144 192.168.226.141 KRBS 1595 TGS-REP
- 0.. 192.168.226.141 192.168.226.144 DCERPC 487 Bind: call_id: 2, Fragment: Single, 3 c
- 0.. 192.168.226.144 192.168.226.141 DCERPC 338 Bind_ack: call_id: 2, Fragment: Single,
- 0.. 192.168.226.141 192.168.226.144 DCERPC 274 Alter_context: call_id: 2, Fragment: S
- 0.. 192.168.226.141 192.168.226.144 LDAP 384 bindRequest(63) "<ROOT>" sasl
- 0.. 192.168.226.144 192.168.226.141 LDAP 264 bindResponse(63) success
- 0.. 192.168.226.141 192.168.226.144 LDAP 378 bindRequest(67) "<ROOT>" sasl
- 0.. 192.168.226.144 192.168.226.141 LDAP 264 bindResponse(67) success
- 55.. 192.168.226.141 192.168.226.144 KRBS 1509 TGS-REQ
- 55.. 192.168.226.144 192.168.226.141 KRBS 1416 TGS-REP
- 55.. 192.168.226.141 192.168.226.144 KRBS 1329 TGS-REQ
- 55.. 192.168.226.144 192.168.226.141 KRBS 1260 TGS-REP
- 55.. 192.168.226.141 192.168.226.144 SMB2 1364 Session Setup Request
- 55.. 192.168.226.144 192.168.226.141 SMB2 314 Session Setup Response
    
```

```

> Record Mark: 1483 bytes
> as-rep
  pvno: 5
  msg-type: krb-as-rep (11)
  > padata: 1 item
    crealm: SAMPLE.LOCAL
  < cname
  > ticket
    tkt-vno: 5
    realm: SAMPLE.LOCAL
    > sname
    > enc-part
      etype: eTYPE-AES256-CTS-HMAC-SHA1-96 (18)
      kvno: 2
      cipher: 3835efb67432dac5c36078427a1dc47ef83070f3677ef7b29874fcabbe2f7462d31efff.
    > enc-part
    
```

Figure 13: TGT issued to the user in AS\_REP

```

- 55.. 192.168.226.141 192.168.226.144 KRBS 1509 TGS-REQ
- 55.. 192.168.226.144 192.168.226.141 KRBS 1416 TGS-REP
- 55.. 192.168.226.141 192.168.226.144 KRBS 1329 TGS-REQ
- 55.. 192.168.226.144 192.168.226.141 KRBS 1260 TGS-REP
- 55.. 192.168.226.141 192.168.226.144 SMB2 1364 Session Setup Request
- 55.. 192.168.226.144 192.168.226.141 SMB2 314 Session Setup Response
    
```

```

> sname
  > enc-part
    etype: eTYPE-ARCFOUR-HMAC-HDS (23)
    kvno: 2
    cipher: dce1ce2f7726a1240a089d4d096a39806d6d1fc5ddac602d5fd7e2baf72cbf177645cbea.
  > authenticator
  > PA-DATA pA-PAC-OPTIONS
    Padding: 0
  > kdc-options: 40810000
    realm: SAMPLE.local
  > sname
    name-type: KRBS-NT-SRV-INST (2)
    > sname-string: 2 items
      SNameString: cifs
      SNameString: MLNK-DC
    till: 2037-09-13 02:48:05 (UTC)
    nonce: 853292803
  > etype: 5 items
  > enc-authorization-data
    
```

Figure 14: TGT presented by the user in TGS\_REQ

Figure 15 illustrates the algorithm employed for automating traffic analysis.

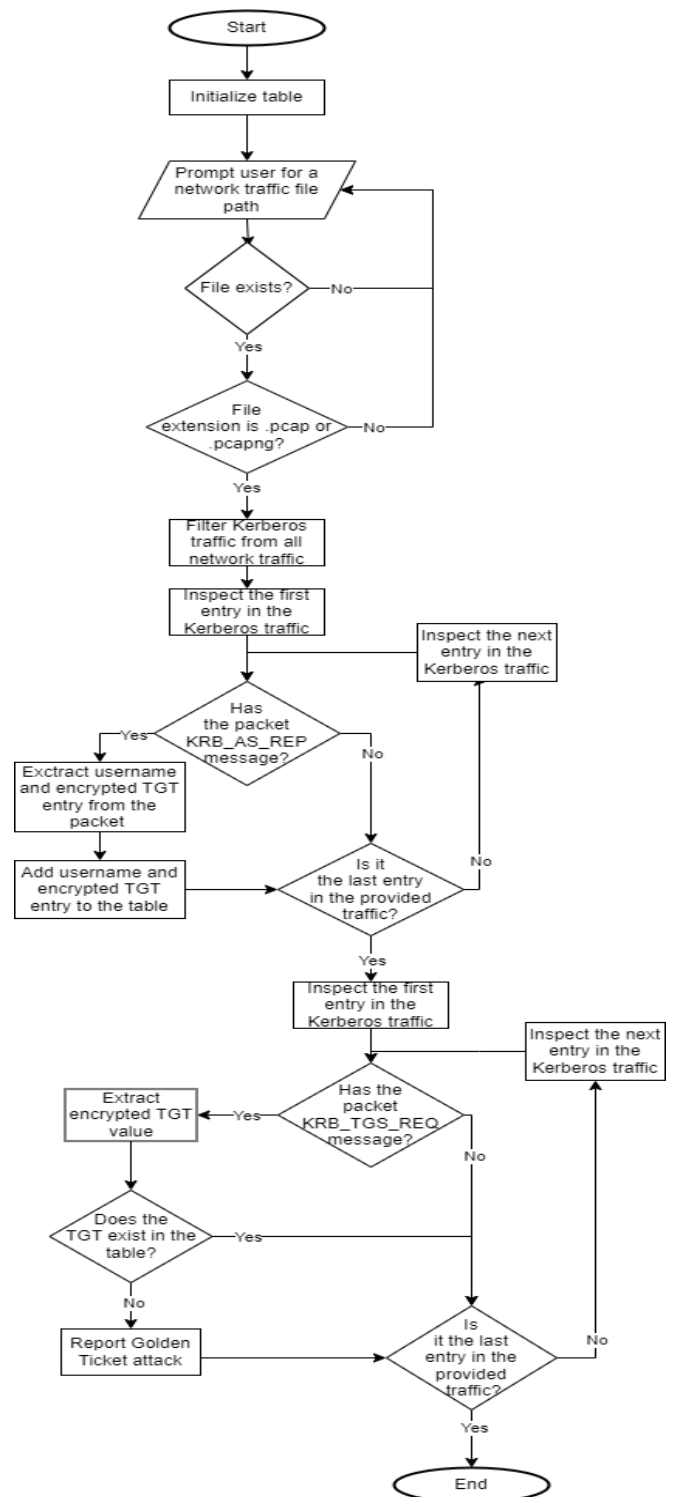


Figure 15: Golden Ticket detection algorithm

The described algorithm was implemented using Python code, and the following screenshot illustrates the algorithm's performance when applied to a Wireshark file containing recorded legitimate activity. The traffic was analyzed, and the TGT ticket issued to the Administrator was recorded in the table.



```
Please provide a Wireshark file to scan for Golden Tickets:
scannet101.onidmco.1.2.pwww
New entry added to table: ['Administrator', '8ab0aba64322743d897c649e3a736abb7e6fa61a5dff7be5b
The whole table now looks like this:
[['Administrator', '8ab0aba64322743d897c649e3a736abb7e6fa61a5dff7be5bcc0779e1d0851607d08e97d51
Process finished with exit code 0
```

Figure 16: Code execution results for legitimate activity

The screenshot provided below displays the algorithm's results for the Golden Ticket activity. Initially, a table was created containing the TGT issued by the DC. Subsequently, the falsified TGT was detected by comparing the user-provided ticket with the one that exists in the table.

```
Please provide a Wireshark file to scan for Golden Tickets:
scannet101.onidmco.1.2.pwww
New entry added to table: ['jdoe', '3835efb67432dac5c36078427a1dcf47ef83070f3677ef7b29874fcabbe
The whole table now looks like this:
[['jdoe', '3835efb67432dac5c36078427a1dcf47ef83070f3677ef7b29874fcabbe2f7462d31efffeb12cd32d65
! GOLDEN TICKET DETECTED !
TGT cipher was not issued by KDC:
dce1ce2f7726a1240a089d4d096a39806d6d1fc5ddac602d5fd7e2baf72cbf177645cbea4e03f8a85adacde44453e5
Source IP: 192.168.226.141
Destination IP: 192.168.226.144
! GOLDEN TICKET DETECTED !
TGT cipher was not issued by KDC:
dce1ce2f7726a1240a089d4d096a39806d6d1fc5ddac602d5fd7e2baf72cbf177645cbea4e03f8a85adacde44453e5
Source IP: 192.168.226.141
Destination IP: 192.168.226.144
Process finished with exit code 0
```

Figure 17: Code execution results for malicious activity

### 3. CONCLUSIONS AND FUTURE WORK

This research paper analyzes the mechanisms of authentication and authorization in network environments based on Active Directory. It explores the operational principles and potential scenarios of Kerberos Golden Ticket attack, which enable attackers to establish domain persistence, escalate their privileges, and access protected confidential resources without proper authorization. The study identifies signatures that can be employed to detect these attacks and offers an algorithm for effective and accurate attack detection. Timely detection is crucial in networks containing sensitive data and information with limited access.

As part of future work, the focus will be on algorithm optimization and adaptation for real-time traffic scanning and analysis. This will enable the practical application of our findings in live network environments, further enhancing security measures and fortifying defenses against evolving threats.

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